

## PRINCIPLES OF COMPILER DESIGN

Time: Three hours

Full Marks: 100

*Different parts of the same question should be answered together.*

CO1 [20]	<p>[1] Answer any one out of (a) and (b) from this block.</p> <p>(a) A lexical analyzer uses the following patterns to recognize three tokens T1, T2, and T3 over the alphabet {a,b,c}.</p> <p>T<sub>1</sub>: a?(b c)*a  T<sub>2</sub>: b?(a c)*b  T<sub>3</sub>: c?(b a)*c</p> <p>(i) Design DFAs for the three tokens.  (ii) Design an NFA for the Lexical Analyzer.  (iii) Convert the NFA produced in part (ii) to a DFA <span style="float: right;">[3x3+1+10]</span></p> <p>(b) Answer the following questions.</p> <p>(i) Design a DFA for C++ style comments using the following <math>\Sigma = \{/, *, a\}</math>.  (ii) Design a Lexical Analyzer based on the DFA produced in (i). <span style="float: right;">[10+10]</span></p>
CO2 [20]	<p>[2] Answer any two out of (a), (b), and (c) from this block.</p> <p>(a) Construct the LL(1) parsing table for the following grammar after suitable modifications</p> <p style="margin-left: 40px;">S → (L)   a  L → L * S   S  <math>\Sigma = \{a, *, (, )\}</math> <span style="float: right;">[10]</span></p> <p>(b) Consider the following grammar G. Construct the LL(1) parsing table.</p> <p style="margin-left: 40px;">S → F   H  F → p   c  H → d   c <span style="float: right;">[10]</span></p> <p>(c) Consider the following grammar</p> <p style="margin-left: 40px;">R → R ∪ R   R.R   R*   (R)   a   b  Construct a LL(1) parse table for this grammar. <span style="float: right;">[10]</span></p>
CO3 [20]	<p>[3] Consider the following grammar.</p> <p style="margin-left: 40px;">E → TT  T → +T   id</p> <p>Construct the LR(1) Sets of Items. There are 10 sets.  Which LR(1) sets can be merged to create LALR sets? <span style="float: right;">[16+4]</span></p>

CO4 [20]	<p>[4] Answer two questions out of (a), (b), and (c) from this block.</p> <p>(a) Consider the following grammar for C style array declaration.</p> <pre>L → id[num]      L [num]</pre> <p>Define the syntax directed translation for generating three-address code for sentences obtained using this grammar. <span style="float: right;">[10]</span></p> <p>(b) Consider a programming language in which programmers define variables in the following manner.</p> <pre>int a, b, c ; float d, e, f ;</pre> <p>Consider another programming language in which programmers define variables in the following manner.</p> <pre>a, b, c : integer; d, e, f : real;</pre> <p>(i) Define a grammar for declarations representing the first style mentioned. Assume <i>id</i>, <i>'</i>, <i>;</i>, <i>int</i>, <i>float</i> to be terminals.</p> <p>(ii) Write a translation scheme to translate the definitions of variables written in the first style to definitions in the form mentioned in second style. <span style="float: right;">[4+6]</span></p> <p>(c) Consider the syntax of the do-loop in FORTRAN as described below. Design a translation scheme to convert it to three address code.</p> <p>The DO statement repeatedly executes a set of statements.</p> <pre>DO s loop-control</pre> <p>where <i>s</i> is a statement number.</p> <p>The form of <i>loop-control</i> is</p> <pre>variable = e1, e2 [, e3]</pre> <p><i>variable</i> - Variable of type integer, real, or double precision.  <i>e1, e2, e3</i> - Integer constants, specifying initial, limit, and increment (or decrement) values respectively.</p> <p>Description of DO loop</p> <p>A DO loop consists of the following:</p> <ul style="list-style-type: none"> <li>• DO statement</li> <li>• Set of assignment statements called a block</li> <li>• Terminal statement, a CONTINUE statement</li> </ul> <p><b>Restrictions</b></p> <ul style="list-style-type: none"> <li>• The DO variable must not be modified in any way within the range of the DO loop. In FORTRAN, every assignment statement appears in a separate line.</li> </ul> <p style="text-align: right;">[10]</p>
CO5 [20]	<p>[5] Answer any two out of (a), (b), and (c) from this block.</p> <p>(a) Convert the following code to a three-address code, identify the basic blocks and draw the control flow graph.</p> <pre>w = 0; x = x+ y; y = 0; if (x &gt; z) {     y = x;     x++; }</pre>

```
else {  
    y = z;  
    z++;  
}  
w = x+z;
```

[3+5+2]

(b) Consider the following code snippet.

```
for (i = 0, i < n; i++) {  
    for (j = 0; j < n; j++) {  
        if (i % 2) {  
            x += (4*j + 5*i);  
            y += (7 + 4*j);  
        }  
    }  
}
```

With respect to the above code snippet, answer the following questions:

- i) What common sub-expression(s) can be eliminated?
- ii) Identify the loop invariant computation?
- iii) Is there any scope of strength reduction? Explain
- iv) What is the possibility of dead code elimination?

[3+2+3+2]

(c) Draw the hierarchical symbol table for the following code segment.

```
void proc()  
{  
    int one_1;  
    int one_2;  
    {  
        int one_3;  
        int one_4;  
    }  
    int one_5;  
    {  
        int one_6;  
        int one_7;  
    }  
}  
struct xyz  
{  
    int a;  
    int b;  
}
```

[10]